

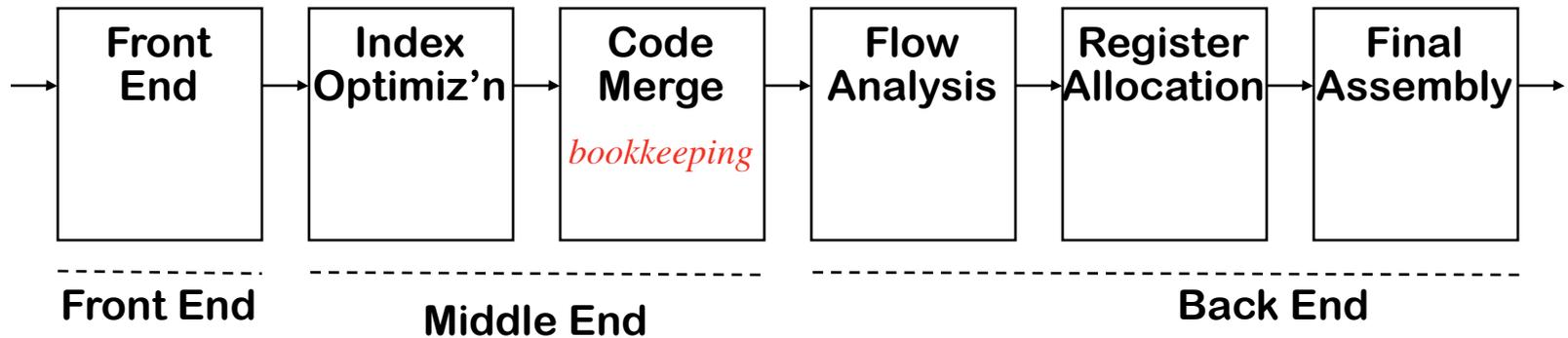


# Historic Compilers

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# Classic Compilers



1957: The FORTRAN Automatic Coding System

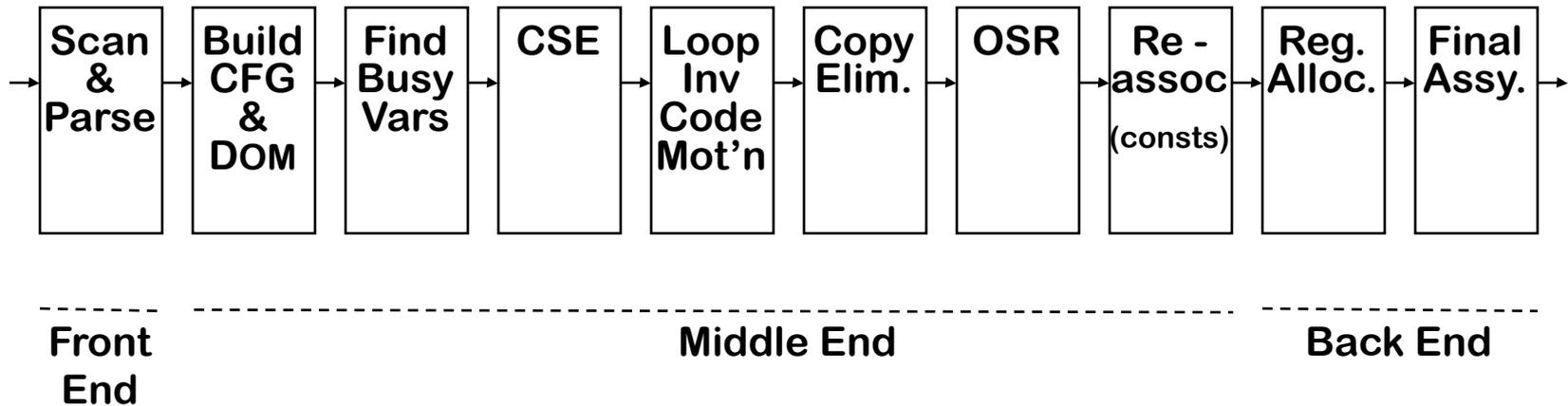
- Six passes in a fixed order
- Generated good code

Assumed unlimited index registers

Code motion out of loops, with ifs and gotos

Did flow analysis & register allocation

# Classic Compilers

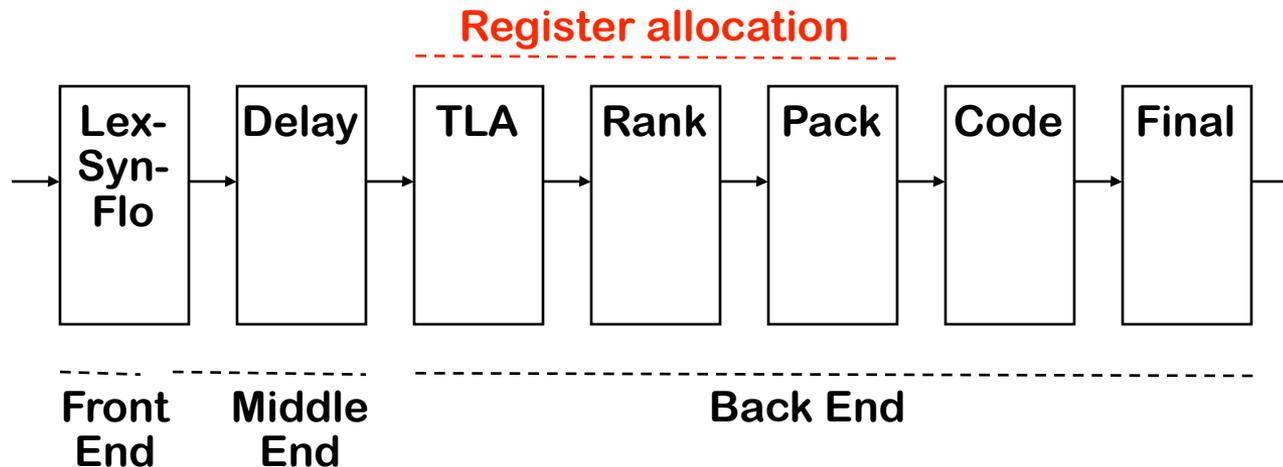


## 1969: IBM's FORTRAN H Compiler

- Used low-level IR (quads), identified loops with dominators
- Focused on optimizing loops ("inside out" order)
  - *Passes are familiar today*
- Simple front end, simple back end for IBM 370



# Classic Compilers



1975: BLISS-11 compiler (Wulf et al., CMU)

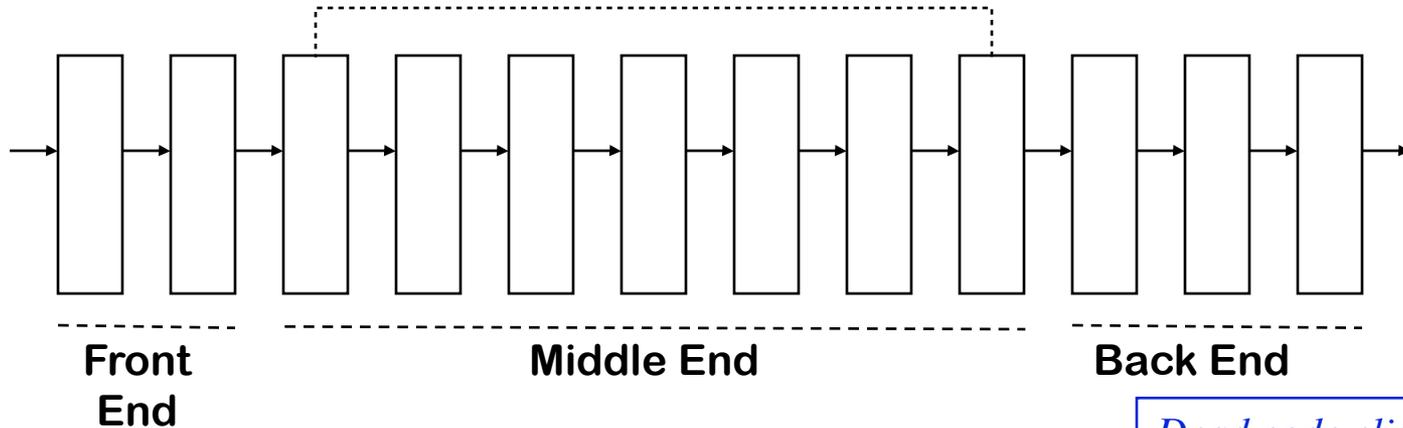
- The great compiler for the PDP-11
- Seven passes in a fixed order
- Focused on code shape & instruction selection

Basis for early VAX & Tartan Labs compilers

LexSynFlo did preliminary flow analysis

Final included a grab-bag of peephole optimizations

# Classic Compilers



## 1980: IBM's PL.8 Compiler

- Many passes, one front end, several back ends
- Collection of 10 or more passes

Repeat some passes and analyses

Represent complex operations at 2 levels

Below machine-level IR

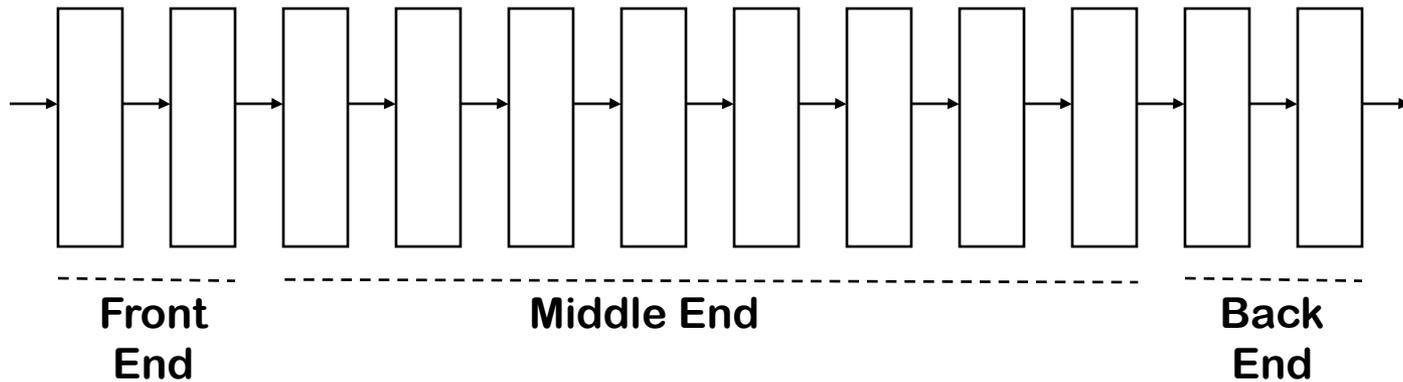
*Multi-level IR  
has become  
common wisdom*

*Dead code elimination  
Global cse  
Code motion  
Constant folding  
Strength reduction  
Value numbering  
Dead store elimination  
Code straightening  
Trap elimination  
Algebraic reassociation*



# Classic Compilers

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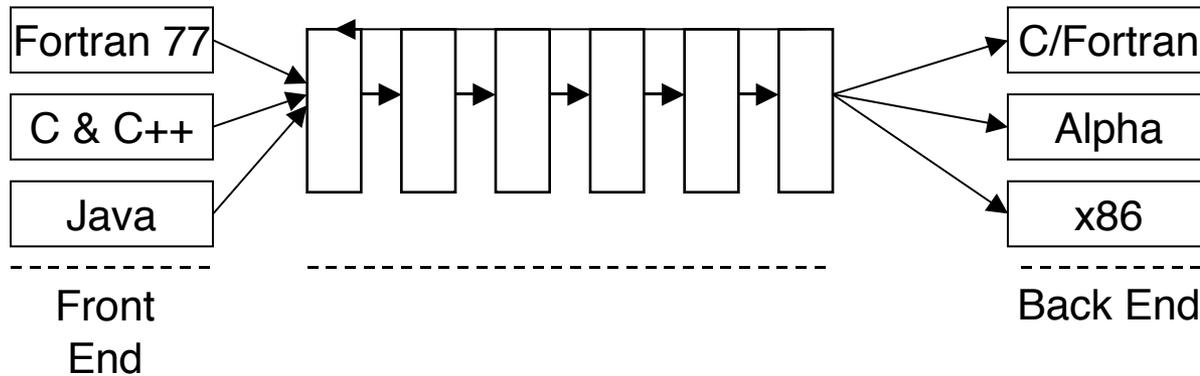
## 1986: HP's PA-RISC Compiler

- Several front ends, an optimizer, and a back end
- Four fixed-order choices for optimization (9 passes)
- Coloring allocator, instruction scheduler, peephole optimizer



# Classic Compilers

## 1999: The SUIF Compiler System



### Another classically-built compiler

- 3 front ends, 3 back ends
- 18 passes, configurable order
- Two-level IR (High SUIF, Low SUIF)
- Intended as research infrastructure

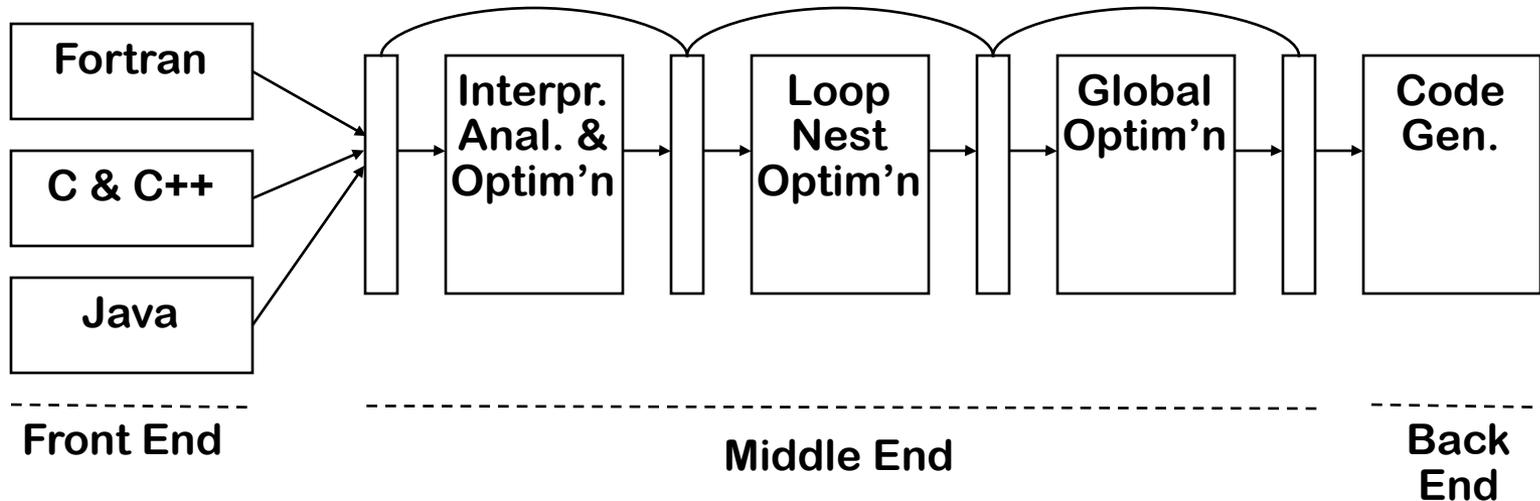
*SSA construction*  
*Dead code elimination*  
*Partial redundancy elimination*  
*Constant propagation*  
*Global value numbering*  
*Strength reduction*  
*Reassociation*  
*Instruction scheduling*  
*Register allocation*

*Data dependence analysis*  
*Scalar & array privatization*  
*Reduction recognition*  
*Pointer analysis*  
*Affine loop transformations*  
*Blocking*  
*Capturing object definitions*  
*Virtual function call elimination*  
*Garbage collection*



# Classic Compilers

2000: The SGI Pro64 Compiler (now Open64 from Intel)

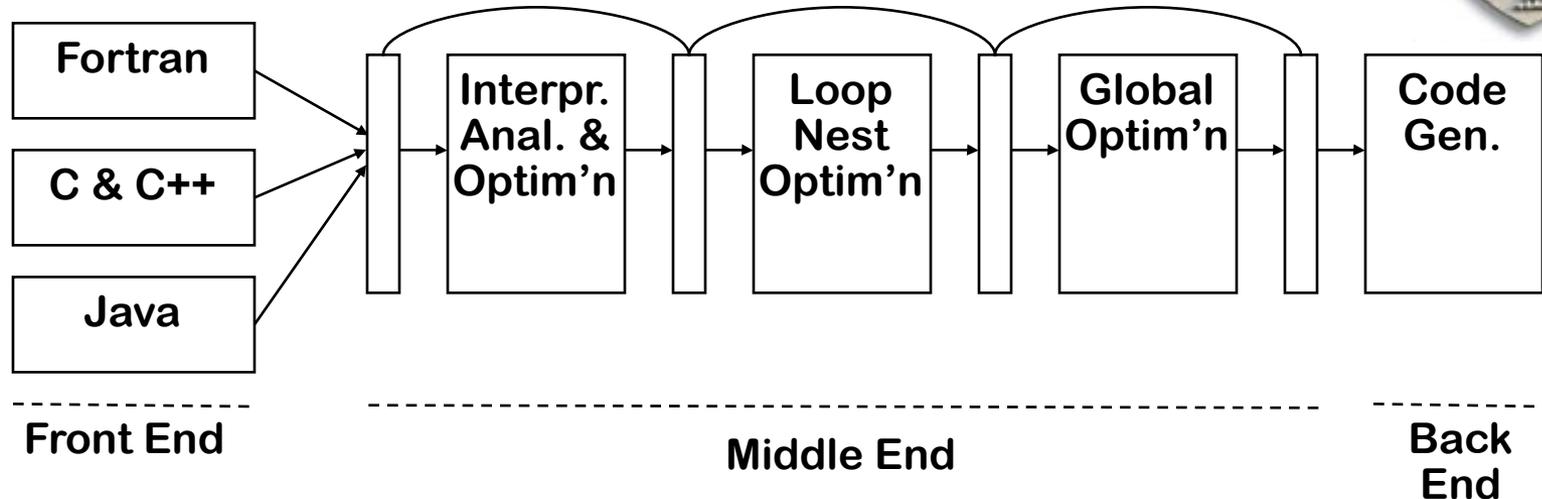


Open source optimizing compiler for IA 64

- 3 front ends, 1 back end
- Five-levels of IR
- Gradual lowering of abstraction level

# Classic Compilers

2000: The SGI Pro64 Compiler (now Open64 from Intel)



## Interprocedural

*Classic analysis*  
*Inlining (user & library code)*  
*Cloning (constants & locality)*  
*Dead function elimination*  
*Dead variable elimination*

## Global Optimization

*SSA-based analysis & opt'n*  
*Constant propagation, PRE,*  
*OSR+LFTR, DVNT, DCE*  
*(also used by other phases)*

## Loop Nest Optimization

*Dependence analysis*  
*Parallelization*  
*Loop transformations (fission, fusion,*  
*interchange, peeling, tiling, unroll & jam)*  
*Array privatization*

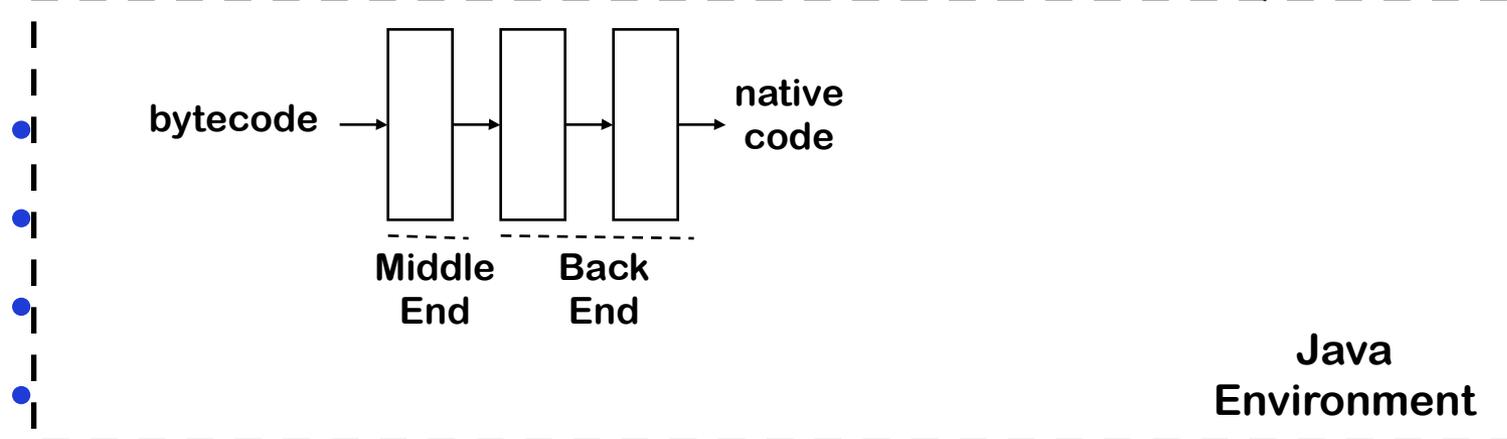
## Code Generation

*If conversion & predication*  
*Code motion*  
*Scheduling (inc. sw pipelining)*  
*Allocation*  
*Peephole optimization*



# Classic Compilers

Even a 2000 JIT fits the mold, albeit with fewer passes



- Front end tasks are handled elsewhere
- Few (if any) optimizations
  - Avoid expensive analysis*
  - Emphasis on generating native code*
  - Compilation must be profitable*